

# Pseudorandom Graphs and the Green-Tao Theorem

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Based on joint work with David Conlon and Jacob Fox

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## A progression of theorems on progressions

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Erdős–Turán conjecture is true.

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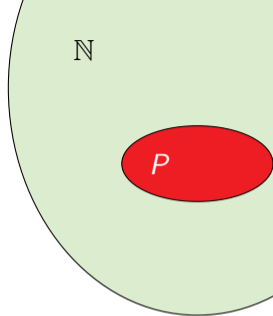
Prime number theorem:  $\frac{\# \text{ primes up to } N}{N} \sim \frac{1}{\log N}$

Our main advance, then, lies not in our understanding of the primes but rather in what we can say about *arithmetic progressions*.

Ben Green  
*Clay Math Proceedings 2007*

# Proof strategy of Green–Tao theorem

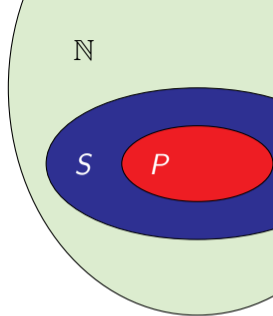
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## Proof strategy of Green–Tao theorem

$P$  = prime numbers,  $S$  = “almost primes”

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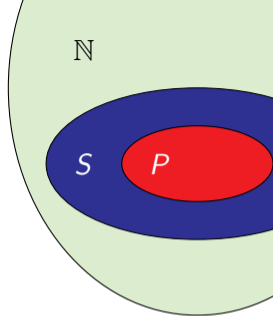
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If  $S \subset \mathbb{N}$  satisfies certain pseudorandomness conditions, then every subset of  $S$  with positive relative density contains long APs.



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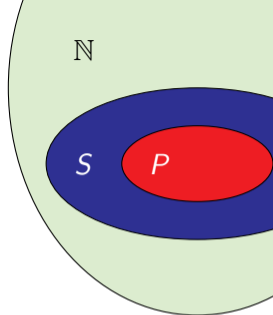
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Step 1:

Relative Szemerédi theorem (informally)

If  $S \subset \mathbb{N}$  satisfies certain pseudorandomness conditions, then every subset of  $S$  with positive relative density contains long APs.

Step 2: Construct a superset of primes that satisfies the pseudorandomness conditions.



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What pseudorandomness conditions?

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Does relative Szemerédi theorem hold with weaker and more natural pseudorandomness hypotheses?



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1. Linear forms condition
  2. Correlation condition ← no longer needed

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Does relative Szemerédi theorem hold with weaker and more natural pseudorandomness hypotheses?

## Theorem (Conlon–Fox–Z. '15)

Yes! A weaker linear forms condition suffices.

# Relative Szemerédi theorem

$k$ -AP-free: contains no  $k$ -term arithmetic progressions

## Szemerédi's theorem (1975)

If  $A \subseteq \mathbb{Z}/N\mathbb{Z}$  is  $k$ -AP-free, then  $|A| = o(N)$ .

## Relative Szemerédi theorem (Conlon–Fox–Z.)

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Earlier versions of relative Roth theorems with other pseudorandomness hypotheses:

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What does it mean for a set to be pseudorandom?

A: It resembles a random set in certain statistics

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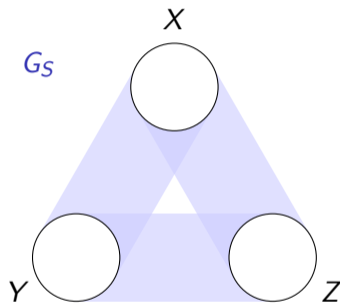
Other ways that graphs can be pseudorandom: eigenvalues, edge discrepancy

Equivalent for dense graphs, but not for sparse graphs

(Thomason '87, Chung–Graham–Wilson '89)

## Graphs and 3-APs (3-term arithmetic progression)

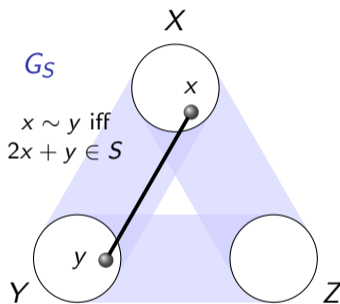
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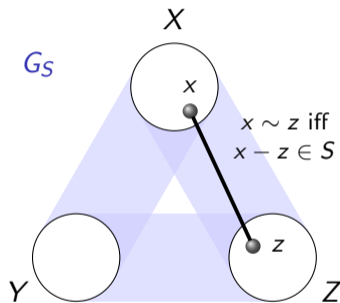
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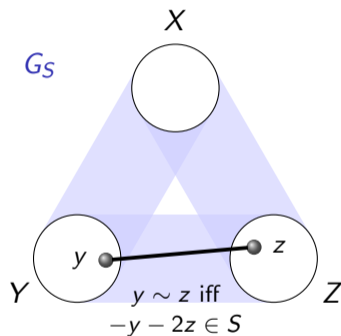
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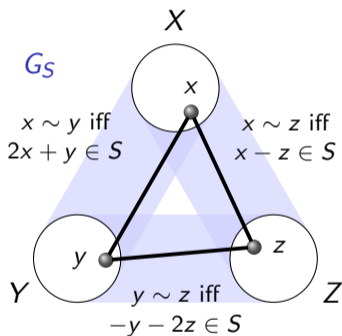
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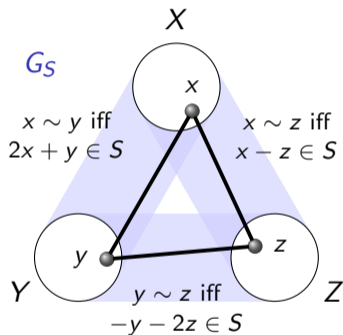


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Triangle  $xyz$  in  $G_S \iff$

$$2x + y, x - z, -y - 2z \in S$$



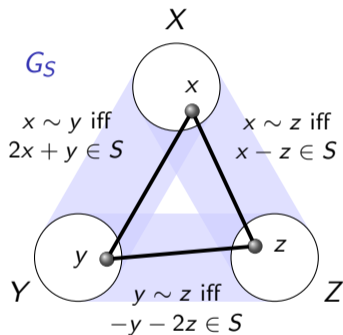
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## Roth's theorem (1952)

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## Relative Roth theorem (Conlon–Fox–Z.)

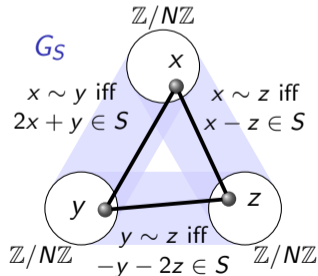
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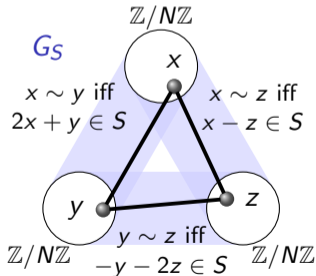


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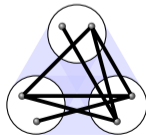
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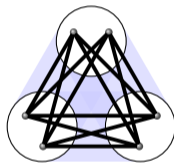
$G_S$  has asymptotically the same  $H$ -density as a random graph for every  $H \subseteq K_{2,2,2}$



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$S \subset \mathbb{Z}/N\mathbb{Z}$  satisfies the **3-linear forms condition** if, for uniformly random  $x_0, x_1, y_0, y_1, z_0, z_1 \in \mathbb{Z}/N\mathbb{Z}$ , the probability that

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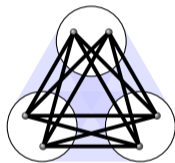


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## Relative Szemerédi theorem (Conlon–Fox–Z.)

Fix  $k \geq 3$ . If  $S \subseteq \mathbb{Z}/N\mathbb{Z}$  satisfies the  $k$ -linear forms condition, and  $A \subseteq S$  is  $k$ -AP-free, then  $|A| = o(|S|)$ .

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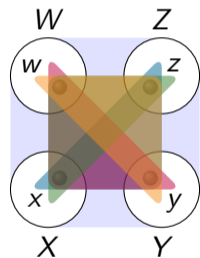
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4-AP with common diff:  $-w - x - y - z$



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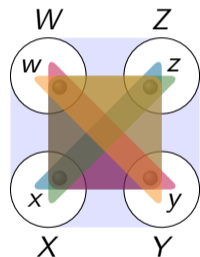
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4-AP with common diff:  $-w - x - y - z$

**4-linear forms condition:** If  $H$  is a subgraph of the 2-blow-up of the tetrahedron, then the  $H$ -density in the above hypergraph is asymptotically same as random

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## Roth's theorem: from one 3-AP to many 3-APs

### Roth's theorem

Let  $\delta > 0$ , every  $A \subset \mathbb{Z}/N\mathbb{Z}$  with  $|A| \geq \delta N$  contains a 3-AP if  $N$  is sufficiently large.

By an averaging argument (Varnavides), we get many 3-APs:

### Roth's theorem (counting version)

Every  $A \subset \mathbb{Z}/N\mathbb{Z}$  with  $|A| \geq \delta N$  contains  $\geq c(\delta)N^2$  many 3-APs for some  $c(\delta) > 0$ .

## Transference

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$\implies$  **relative Roth theorem** (also works for  $k$ -AP)

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$\implies$  relative Roth theorem (also works for  $k$ -AP)

## Dense model

What does it mean for

$$\text{(dense)} \quad \tilde{A} \subset \mathbb{Z}/N\mathbb{Z}$$

to be a good approximation (dense model) of

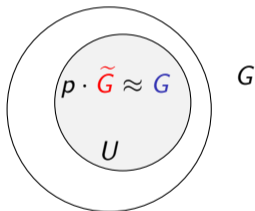
$$\text{(sparse)} \quad A \subset S \subset \mathbb{Z}/N\mathbb{Z} ?$$

## Dense model

Let  $\tilde{G}$  (dense) and  $G$  (sparse) be two graphs on the same set of  $N$  vertices

We say that  $\tilde{G}$  is a good  $p$ -dense model of  $G$  if  $p \cdot \tilde{G} \approx G$  in terms of the number of edges when restricted to every vertex subset, i.e.,

$$|p \cdot e_{\tilde{G}}(U) - e_G(U)| = o(pN^2) \quad \forall U \subset V(G) = V(\tilde{G})$$



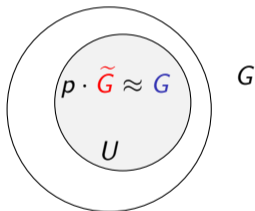


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We say that  $\tilde{A} \subset \mathbb{Z}/N\mathbb{Z}$  is a good  $p$ -dense model of  $A \subset \mathbb{Z}/N\mathbb{Z}$  if

$\text{CayleySumGraph}(\mathbb{Z}/N\mathbb{Z}, \tilde{A})$  is a good  $p$ -dense model of  $\text{CayleySumGraph}(\mathbb{Z}/N\mathbb{Z}, A)$

$\text{CayleySumGraph}(G, A)$  has vertex set  $G$ , and  $x \sim y$  iff  $x + y \in A$

## Dense model theorem

If  $\mathbb{Z}/N\mathbb{Z}$  is a good  $p$ -dense model of  $S \subset \mathbb{Z}/N\mathbb{Z}$  with  $p = |S|/N$ , then every  $A \subset S$  has a good  $p$ -dense model  $\tilde{A} \subset \mathbb{Z}/N\mathbb{Z}$ .

Proof ideas: Hahn–Banach theorem/linear programming duality

Originally Green–Tao and Tao–Ziegler. Simplified by Gowers and Reingold–Trevisan–Tulsiani–Vadhan. Specialized to this form in Z.

# Transference

Let  $S \subset \mathbb{Z}/N\mathbb{Z}$  be pseudorandom with density  $p$ , and

$$\text{(sparse)} \quad A \subset S \subset \mathbb{Z}/N\mathbb{Z}, \quad |A| \geq \delta |S|$$

**Dense model theorem:** One can find a good  $p$ -dense model  $\tilde{A}$  of  $A$ :

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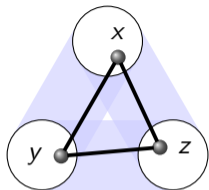
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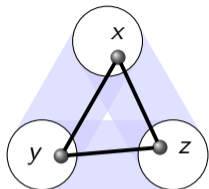


### Triangle counting lemma, dense setting

Let  $G$  and  $\tilde{G}$  be (tripartite) graphs on the same vertex set, such that  $\tilde{G}$  is a good 1-dense model of  $G$ . Then

$$\text{triangle-density}(G) = \text{triangle-density}(\tilde{G}) + o(1)$$

## Counting lemma



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### Triangle counting lemma, sparse setting (Conlon–Fox–Z.)

(Sparse)  $G \subset \Gamma$  and (dense)  $\tilde{G}$  are (tripartite) graphs on the same vertex set. Suppose

- ▶ “Sparse pseudorandom host graph”  $\Gamma$  has edge density  $p$  and satisfies the 3-linear forms condition (densities of  $H \subset K_{2,2,2}$  are close to random)
- ▶  $\tilde{G}$  is a good  $p$ -dense model of  $G$

Then

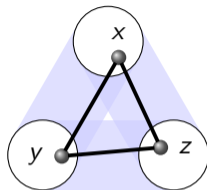
$$\text{triangle-density}(G) = p^3(\text{triangle-density}(\tilde{G}) + o(1))$$

# Counting lemma

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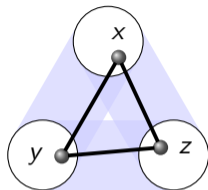


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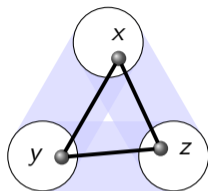


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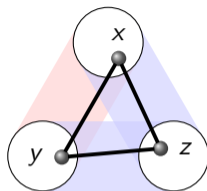
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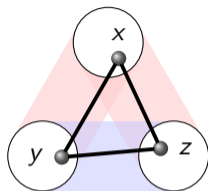
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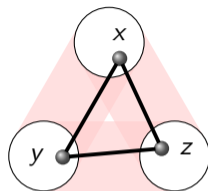
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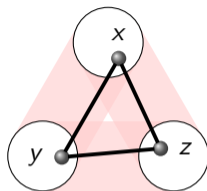
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Fails in the sparse setting (need  $o(p^3)$  error)

# Sparse counting lemma

## Triangle counting lemma, sparse setting (Conlon–Fox–Z.)

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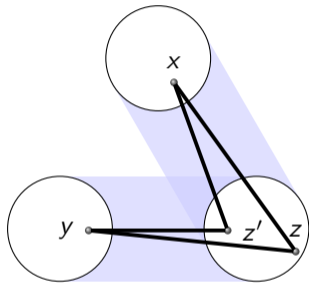
- ▶ “Sparse pseudorandom host graph”  $\Gamma$  has edge density  $p$  and satisfies the 3-linear forms condition (densities of  $H \subset K_{2,2,2}$  are close to random)
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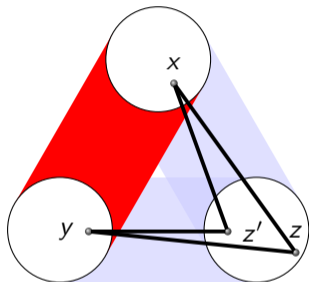
Key new proof ingredient: **densification**

## Densification



$$\mathbb{E}[G(x, z)G(y, z)G(x, z')G(y, z')]$$

# Densification



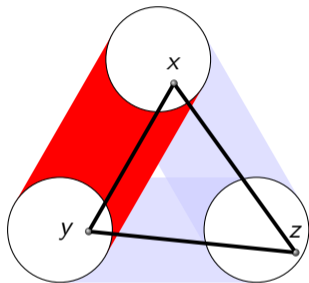
$$\begin{aligned}\mathbb{E}[G(x, z)G(y, z)G(x, z')G(y, z')] \\ = \mathbb{E}[G'(x, y)G(x, z)G(y, z)]\end{aligned}$$

Set  $G'(x, y) := \text{codeg}_G(x, y) / |Z|$

$G'(x, y) = O(p^2)$  for almost all pairs  $(x, y)$ ,  
and thus behaves like a dense weighted graph  
after scaling



## Densification



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Densified  $G(X, Y)$ . Now repeat for  $G(X, Z)$  and  $G(Y, Z)$ .  
Reduce to dense setting.

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$\implies$  relative Roth theorem (also works for  $k$ -AP)

# Relative Szemerédi theorem

## Szemerédi's theorem (1975)

If  $A \subseteq \mathbb{Z}/N\mathbb{Z}$  is  $k$ -AP-free, then  $|A| = o(N)$ .

## Relative Szemerédi theorem (Conlon–Fox–Z.)

If  $S \subseteq \mathbb{Z}/N\mathbb{Z}$  satisfies the  $k$ -linear forms condition, and  $A \subseteq S$  is  $k$ -AP-free, then  $|A| = o(|S|)$ .

## Green–Tao theorem

Every subset of the primes with positive relative density contains arbitrarily long APs.

# Polynomial progressions in the primes

## Polynomial Szemerédi theorem (Bergelson–Leibman 1996)

Every subset of  $\mathbb{N}$  with positive density contains arbitrary **polynomial progressions**, i.e., for every  $P_1, \dots, P_k \in \mathbb{Z}[X]$  with  $P_1(0) = \dots = P_k(0) = 0$ , the subset contains  $x + P_1(y), \dots, x + P_k(y)$  for some  $x$  and  $y > 0$ .

## Polynomial Szemerédi theorem in the primes (Tao–Ziegler 2008)

Every subset of the primes with positive relative density contains arbitrary polynomial progressions.

Using the **densification method**, Tao and Ziegler recently strengthened their result:

- ▶ (2015) existence of *narrow* progressions with polylogarithmic gaps
- ▶ (2018) asymptotics for the number of polynomial patterns in the primes

## Some open problems

- ▶ Can the pseudorandomness hypotheses be further weakened?
- ▶ A **multidimensional relative Szemerédi theorem?**  
Linear forms conditions on  $S \subset \mathbb{Z}/N\mathbb{Z}$  so that every relatively dense  $A \subset S \times S$  contains a  $k \times k$  square grid



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THANK YOU!